

DOGS 28/05

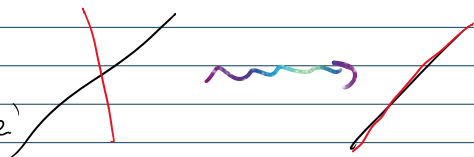
What the hell is derived intersection theory?

I don't know :-

Intersection Th^y

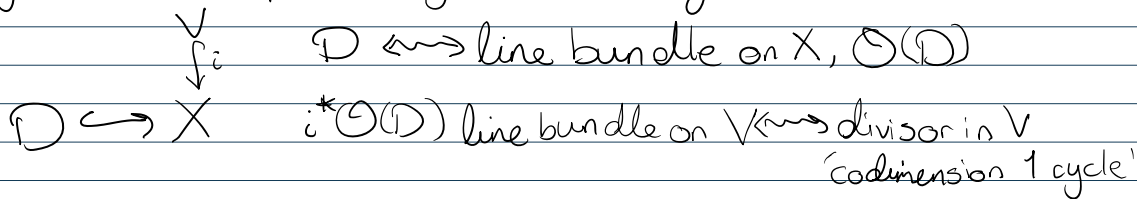
Want sensible numerics in enumeration problems following a Principle of Continuity

Solⁿ Consider 'cycles' upto 'rational equivalence'



'Weil divisors' 'linear equivalence'

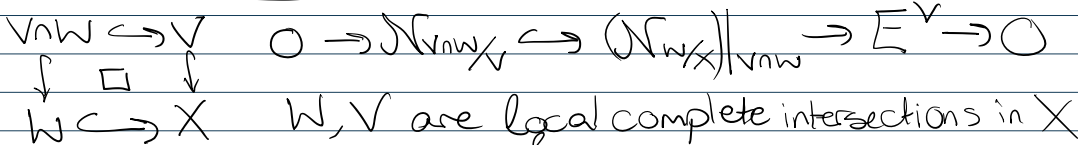
Define intersection product of divisors as follows



Weird to me that $\mathcal{D} \cdot \mathcal{D}$ is somehow codimension 1 in the divisor

Can generalise to higher codim

Excess Intersection Formula



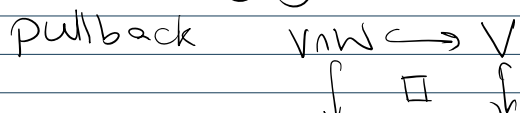
$W \cap V$ is a lci in V

$$V \cdot W = c_{d'}(E^V) \cap [V \cap W]$$

$$d = \text{codim}(W, X), d' = \text{codim}(V \cap W, V)$$

Derived

In any category an intersection of subobjects is given by a

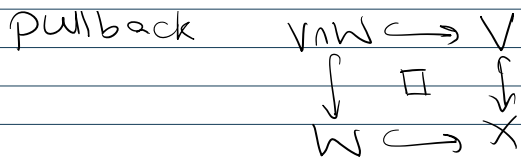


Intersection Theory



Derived Categories

Homological Algebra



Topology, $V \cap W = \text{set-theoretic intersection}$

Scheme theory $(V, \mathcal{O}_V) \cap (W, \mathcal{O}_W) = (V \cap W, \mathcal{O}_V \otimes_{\mathcal{O}_X} \mathcal{O}_W)$

Derived scheme theory $(V, \mathcal{O}_V) \cap (W, \mathcal{O}_W) = (V \cap W, \mathcal{O}_V \otimes^L_{\mathcal{O}_X} \mathcal{O}_W)$

↑ Intersection structure.

$$\mathcal{O}_V \otimes^L \mathcal{O}_W \cong P_V^\bullet \otimes \mathcal{O}_W \cong \mathcal{O}_V \otimes P_W^\bullet \cong P_V^\bullet \otimes P_W^\bullet \rightarrow D(X)$$

Here P_V, P_W are flat resolutions of $\mathcal{O}_V, \mathcal{O}_W$

In the case of $D \subset X$ an effective Cartier divisor

Scheme theory

Derived

$$\begin{aligned}
 (D, \mathcal{O}_D) \cap (D, \mathcal{O}_D) &\cong (D, \mathcal{O}_D \otimes \mathcal{O}_D) \\
 &\cong (D, \mathcal{O}_D)
 \end{aligned}$$

$$\begin{aligned}
 (D, \mathcal{O}_D \otimes^L \mathcal{O}_D) \\
 \cong (D, \{ \mathcal{O}_D(-1) \rightarrow \mathcal{O}_D \})
 \end{aligned}$$

↑ measures excess
Cox

↑ classical

Comparison

$$\text{Tor}_q^{\mathcal{O}_X}(\mathcal{O}_V, \mathcal{O}_W) = H^{-q}(\mathcal{O}_V \otimes^L \mathcal{O}_W)$$

Th^m (Serre)

X regular variety, $V, W \subseteq X$ complementary dimension

s.t. $\dim(X \cap V) = 0$, Then

$$\underbrace{\underline{i}(P, V, W; X)}_{\text{Can define via Fulton}} = \sum_{i=0}^{\infty} (-1)^i \text{length}_{\mathcal{O}_{X,P}} \text{Tor}_i^{\mathcal{O}_{X,P}}(\mathcal{O}_{V,P}, \mathcal{O}_{W,P})$$

More generally;

cycles/rat. eq.

$$\exists \text{ GRR iso } \tau: K_0(X)_{\mathbb{Q}} \xrightarrow{\sim} A_*(X)_{\mathbb{Q}}$$

For any $V, W \subseteq X$, define $\text{Tor}^X(V, W) = \sum (-1)^i [\text{Tor}_i^X(\mathcal{O}_V, \mathcal{O}_W)] \in K_0(X)$

Then $\tau(\text{Tor}^X(V, W)) = V \cdot W + \text{lower dim terms.}$

Excess Intersection Formula [Sca/a '15]

Excess Intersection Formula [Scala '15]

Th^m Let $Y_1, \dots, Y_n \subseteq X$ be lci subvarieties of X nonsing. $/k = \bar{k}$ char $k = 0$.
Assume $\bigcap Y_i = Z$ is lci in X . Define E by

$$0 \rightarrow E \rightarrow (\bigoplus C_{Y_i/X})|_Z \rightarrow C_{Z/X} \rightarrow 0$$

\uparrow
 Excess conormal
 directions

Then $\mathrm{Tor}_q^{\mathcal{O}_X}(\mathcal{O}_{Y_1, \dots, Y_n}) \cong \wedge^q E$.

$$\mathrm{Tor}_q(\mathcal{O}_{Y_1, \dots, Y_n}) \cong H^{-q}(\mathcal{O}_{Y_1} \otimes^L \dots \otimes^L \mathcal{O}_{Y_n})$$

Special case is self-intersection formula

$$\mathrm{Tor}_q^{\mathcal{O}_Y}(\mathcal{O}_Y, \mathcal{O}_Y) \cong \wedge^q C_{Y/X}$$

$$\cong H^{-q}(\mathbb{P}_Y^* \otimes \dots \otimes \mathbb{P}_Y^*)$$

$\mathrm{Tor}_q^{\mathcal{O}_Y}(\mathcal{O}_Y, \mathcal{O}_Y)$ is only really defined up to isomorphism

Need to pick flat models $P_V \rightarrow \mathcal{O}_V, P_W \rightarrow \mathcal{O}_W$ then $\mathrm{Tor}_q^{\mathcal{O}_Y}(\mathcal{O}_Y, \mathcal{O}_Y) := H^{-q}(P_V^* \otimes P_W^*)$

Picking different flat models gives rise to an isomorphism

$$H^{-q}(P_V^* \otimes P_W^*) \rightarrow H^{-q}(Q_V^* \otimes Q_W^*)$$

Local complete intersections = Good because they locally have bdd free resolutions

$\mathrm{Spec}(A), \mathbb{P}_Y$ defined by regular sequence $f_1, \dots, f_r = \mathrm{codim}(Y, X)$

$$K(f_1, \dots, f_r) \text{ defined by } K^{-q}(f_1, \dots, f_r) = \wedge^q A^r \cong A^{\binom{r}{q}}$$

$$d(a_1 \wedge \dots \wedge a_q) = \sum_i (-1)^i f_i a_1 \wedge \dots \wedge \hat{a}_i \wedge \dots \wedge a_r$$

$$K(f_1) \quad A \xrightarrow{f_1} A$$

$$K(f_1, f_2, f_3)$$

$$K(f_1, f_2) = \begin{array}{c} A \xrightarrow{f_2} A \oplus A \xrightarrow{f_1} A \\ \uparrow \quad \uparrow \quad \downarrow \\ A \xrightarrow{f_2} A \xrightarrow{f_2} A \\ \downarrow \quad \downarrow \\ A \xrightarrow{-f_1} A \end{array}$$

$$\begin{array}{c} A \xrightarrow{f_3} A \xrightarrow{f_2} A \xrightarrow{f_1} A \\ \uparrow \quad \uparrow \quad \downarrow \\ A \xrightarrow{f_3} A \xrightarrow{f_2} A \xrightarrow{f_2} A \\ \uparrow \quad \uparrow \quad \downarrow \\ A \xrightarrow{f_1} A \xrightarrow{f_3} A \xrightarrow{-f_2} A \end{array}$$

Sequence f_1, \dots, f_r regular $\Leftrightarrow K(f_1, \dots, f_r) \rightarrow A/(f_1, \dots, f_r)$ a resolution.

Computation

V, W lci $\subseteq X$ nonsing. $/k = \bar{k}$

$$v, v \in \mathcal{K} = \Lambda \text{ mapping } \mathcal{K} \rightarrow \mathcal{K} \rightarrow \mathcal{K}$$

Pick global flat res $P_v \rightarrow \mathcal{O}_v, P_w \rightarrow \mathcal{O}_w$.

$$\text{Local Koszul models } K(f_1, \dots, f_r) \rightarrow \mathcal{O}_v|_U \\ K(g_1, \dots, g_s) \rightarrow \mathcal{O}_w|_U$$

Freeness implies unique lifts (up to homotopy)

$$\begin{array}{ccc} & \rightarrow P_v|_U & P_w|_U \leftarrow \\ & \downarrow & \downarrow \\ K(f_1, \dots, f_r) \rightarrow & \mathcal{O}_v|_U & \mathcal{O}_w|_U \leftarrow K(g_1, \dots, g_s) \end{array}$$

induces a quasi-isomorphism

$$K(f_1, \dots, f_r) \otimes K(g_1, \dots, g_s) \rightarrow P_v|_U \otimes P_w|_U \text{ because of } q\text{-flatness}$$

Hence isomorphisms

$$\varphi: H^{-q}(K(f_1, \dots, f_r, g_1, \dots, g_s)) \rightarrow \text{Tor}_q(\mathcal{O}_v, \mathcal{O}_w)|_U$$

What if we pick different generators?

We get a 'change of basis' map $K: K(f_1, \dots, f_r, g_1, \dots, g_s) \rightarrow K(f'_1, \dots, f'_r, g'_1, \dots, g'_s)$

s.t.

$$\begin{array}{ccc} & \text{Tor}_q(\mathcal{O}_v, \mathcal{O}_w)|_U & \\ \varphi \nearrow & \uparrow & \searrow \varphi' \\ H^{-q}(K(\underline{f}, \underline{g})) & \xrightarrow{H^{-q}K} & H^{-q}(K(\underline{f}', \underline{g}')) \end{array}$$

$$\begin{array}{ccc} & \uparrow & \\ & \mathcal{F}|_U & \\ & \downarrow & \end{array}$$

$$\Rightarrow \text{Tor}_q(\mathcal{O}_v, \mathcal{O}_w) \cong \mathcal{F}$$

Sounds easy and boring

Boring ✓

Self-intersection

$$\text{Isomorphisms } H^{-q}(K(f_1, \dots, f_r) \otimes \Lambda / (f_1, \dots, f_r)) \cong K^{-q}(f_1, \dots, f_r) \otimes \Lambda / (f_1, \dots, f_r) \\ \cong \Lambda^q C_{Y/X}|_U$$

+ gluing data

This is not hard

Same philosophy is used to prove EIF but much harder to show gluing because the local isomorphisms are more involved.

same philosophy is used to prove $\leq + 1$ but much harder to show gluing because the local isomorphisms are more involved.

Reduce to the diagonal

$$\mathcal{O}_Y \otimes^L \dots \otimes^L \mathcal{O}_{Y_n} \cong \Delta^* \mathcal{O}_{Y_1 \times \dots \times Y_n}$$

$$\begin{array}{ccc} \Pi Y_i & \xrightarrow{k} & X \\ w \downarrow & \square & \downarrow \Delta \\ Y_1 \times \dots \times Y_n & \xrightarrow{i_{1 \times \dots \times 1n}} & X^{1n} \end{array}$$

$$\begin{array}{ccc} \Delta^* \mathcal{O}_{Y_1 \times \dots \times Y_n} & \xrightarrow{\eta_k} & k_* k^* \Delta^* \mathcal{O}_{Y_1 \times \dots \times Y_n} & \xrightarrow{L} & i_* \mathcal{O}_Y \otimes^L i_* \mathcal{O}_Y \\ & \cong & k_* w^* (i_{1 \times \dots \times 1n})^* (i_{1 \times \dots \times 1n})_* \mathcal{O}_{Y_1 \times \dots \times Y_n} & \cong & Li^* i_* \mathcal{O}_Y \end{array}$$

self-intersection of the lci
 $Y_1 \times \dots \times Y_n \hookrightarrow X^{1n}$

Global diagram

$$\begin{array}{ccc} \text{Tor}_q^{\mathbb{Z}}(\mathcal{O}_{Y_1}, \dots, \mathcal{O}_{Y_n}) & \xrightarrow{H^{-q} \eta_k} & H^{-q}(k_* k^* \Delta^* \mathcal{O}_{Y_1 \times \dots \times Y_n}) \\ & & \downarrow \cong \varphi \\ \Lambda^q E & \hookrightarrow & \Lambda^q(\bigoplus \mathcal{O}_{Y_i, X})|_W \end{array}$$

Just need to check that

- 1) $H^{-q} \eta_k$ is injective
- 2) Its image agrees with the image of $\Lambda^q E$

These checks are local. \therefore no gluing required